

Comparative Analysis of Vedic Multiplier by Using Different Adder Logic Style at Deep Submicron Technology

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Abstract- In the recent years growth of the portable electronic is forcing the designers to optimize the existing design for better performance. Multiplication is the most commonly used arithmetic operation in various applications like DSP processor, math processor and in various scientific arithmetic circuits. Overall performance of the VLSI system is strongly depends on the performance of arithmetic circuits like multiplier. Different types of multiplier are available in the literature depending on the requirement. Designer used different type of multiplier. Power consumption in a multiplier is more prominent design factor in addition with the performance, because portable equipment needs larger battery backup which is only possible through low power design. This work focuses on the reduction of the power dissipations which is showing an ever-increasing growth with the scaling down of the technologies. The power consumption is an important issue that has lead to multiplier designer to adopt different technique to reduce power dissipations. In this thesis work, a Vedic algorithm is used to reduce the power consumption in multiplier used. This work focus on the analysis of power dissipation, delay, area of Vedic Multiplier using different adder logic style at 180nm technology by using tanner EDA tool.

I. INTRODUCTION

Digital multiplication is one of the most basic functions in arithmetic circuits. The ubiquity of this operation in computing has given rise to a large number of multiplier implementations, each with different specifications and goals. Some applications require wide dynamic range, others need high precision, while in some cases, neither of these characteristic is very tightly specified. Digital multiplication is used as

Opposed to analog when high precision is an issue; it is fairly straightforward to make digital

Multipliers are accurate as the application requires. Precision required for multiplication

Varies by function. The basic operation in these designs is integer multiplication. In floating

Point multipliers, integer multiplication units are sub-blocks of the greater floating point unit. Signed versus unsigned techniques have an impact on the design, and some other techniques have been suggested for manipulating the bit representation of

numbers to generate power savings. However, the primary consideration in multipliers has been and continues to be delay. The demand for high speed processing has been increasing as a result of expanding computer and signal processing applications. Higher throughput arithmetic operations are important to achieve the desired performance in many real-time signal and image processing applications [1]. One of the key arithmetic operations in such applications is multiplication and the development of fast multiplier circuit has been a subject of interest over decades. Reducing the time delay and power consumption are very essential requirements for many applications. This work presents Vedic multiplier with different circuit style of adder. Multiplier based on Vedic Mathematics is one of the fast and low power multiplier.

Need for Vedic Mathematics:-

- It consumes the less time.
- It reduces the scratch from finger.
- It is magical tool for the calculations.
- It provides the one line answer.
- It gives the more concentrations to works.

II. PRESENT WORK

Vedic Mathematics: - The proposed Vedic multiplier is based on the Vedic multiplications formulae (Sutras). This Sutra has been traditionally used for the multiplications of two numbers in the decimal number system. In this work we apply the same ideas to the binary number system to make the proposed algorithm compatible with the digital hardware. Vedic multiplications based on some algorithm, one of them is discussed below:

Urdhava Triyakbhyam :- This multiplier is based on an algorithm Urdhava Triyakbhyam (Vertical & Crosswise) of ancient Indian Vedic multiplications. Urdhava Triyakbhyam Sutras is a general multiplications formula applicable to all cases of multiplication. It literally means "Vertically and crosswise". It is based on a novel technique through which the generations of all partial products can be

done with the concurrent additions of these partial products. The parallelism in generations of partial products and their summations is obtained using Urdhava Triyakbhyam in fig 1. Since the partial products and their sum are calculated in parallel, the multiplier is independent of the clock frequency of the processor. The net advantage is that it reduces the need of microprocessor to operate at increasingly high clock frequency. While a higher clock frequency generally results in increased processing power, its disadvantage is that it also increases power dissipations which results in higher device operating temperature. By adopting the Vedic multiplier, microprocessors designer can easily circumvent these problem to avoid catastrophic device failures.

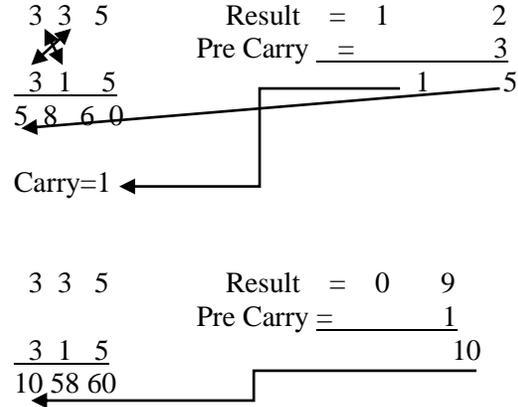
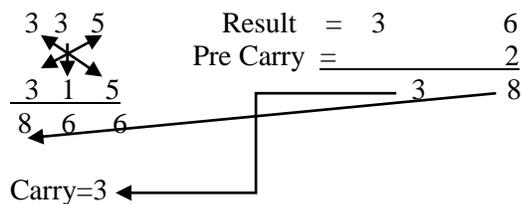
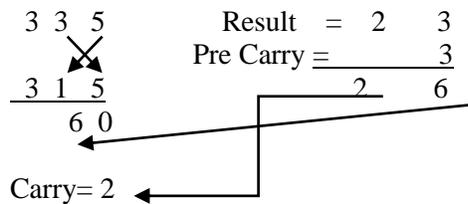
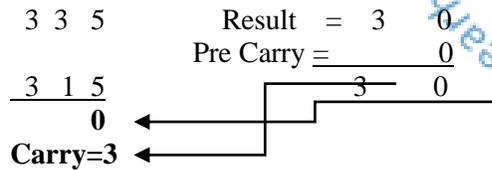


Fig.1:- Multiplications of two decimal numbers by Urdhva Tiryakbhyam

Multiplications of two decimal numbers (335 * 315)

Let us consider the multiplications of two decimal numbers (335 * 315). The digits on the both sides of the lines are multiplied and added with the carry from the previous step. This generates one of the bits of the results and a carry, this carry is added in the next steps and hence a process is going on. If more than one line are added to the previous carry. In each steps least significant bits acts as the results bit and all other bit acts as carry for the next step. Initially the carry is taken to be zero.



Algorithm for N x N bit Vedic multiplier Using Urdhva Tiryakbhyam for two binary numbers.

For NxN multiplication, divided the multiplicand and multiplier into two parts, consisting of (N to N/2-1) bits and (N/2) bits. For example for multiplications of X and B, of 16 bits each, if X=0000001010101101 then its parts will be 00000010 and 10101101. Represent the above mentioned parts of X as Xm and Xl. Similarly for the B, it is Bm and Bl. Now represent X and B as XmXl and BmBl. For A x B we can have

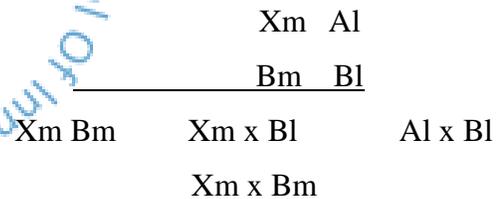


Figure 2:- General Representations of Vedic multiplications

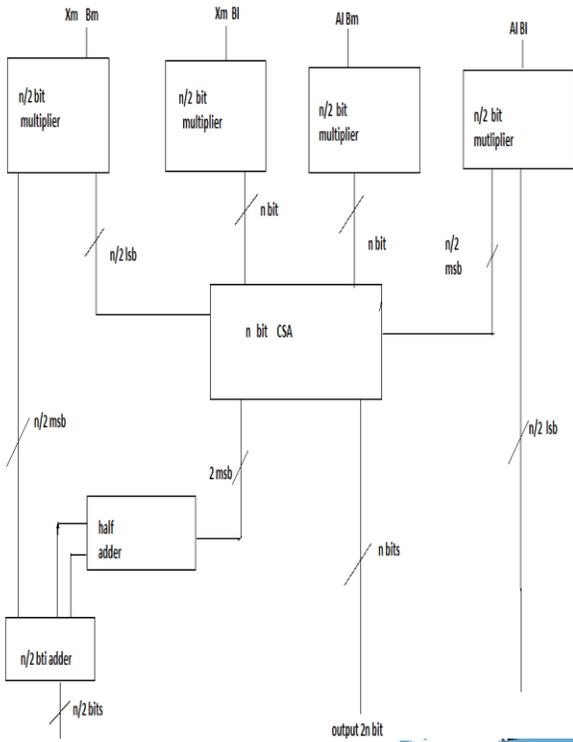


Fig 3:- Complete block diagram for n x n bit Vedic multiplier

Table 1: Truth Table of AND Gate

A	B	Output
0	0	0
0	1	0
1	0	0
1	1	1

III. MULTIPLIER DESIGN

Schematics of different blocks of multiplier

design: - The schematics of different blocks of multiplier design are designed on the TANNER tool by using 180um technology.

Schematic of AND Gate:- In multiplier design, first of all product terms will be generated using AND gates and then these product terms will be added using adders to get the result as explained in earlier chapters. Product term is an AND operation between A bit of one number and the B bit of a second number. If both the inputs to the AND gate are HIGH (1), then a HIGH output (1) results. A LOW output results, if neither or only one input to the AND gate is HIGH. In another sense, the function of AND effectively finds the minimum between two binary digits, just as the OR function finds the maximum. Thus, the output is always 0 except when all the inputs are 1s. The operational truth table of an AND gate is shown in table 4.1 where A, B are the inputs of AND gate and while its schematic is shown in figure 4.3.

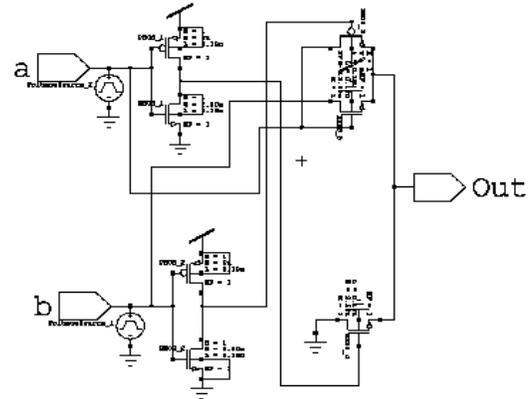


Figure 4:- Schematic of AND Gate

Schematic of Full adder

Basically the Full Adder is type of combinational circuit. The Full Adder are used to overcome the disadvantage of half adder. It will perform the full additions of numbers. The addition of the partial product terms is done with full adder which are generated by AND Gate. The full adder adds binary numbers. A one-bit full adder adds three one-bit numbers, often written as A, B, and C_{in} ; A and B are the operands, and C_{in} is a bit carried in from the next less significant stage [23]. Usually, the full-adder is a component in a cascade of adders, which add 8, 16, 32, etc. binary numbers. A two-bit output is produced

by the circuit. A full adder can be constructed from two half adders by connecting A and B to the input of one half adder, connecting the sum from that to an input to the second adder, connecting C_{in} to the other input and OR the two carry outputs [22]. The operational truth table of Full adder is shown in table 4.2 where A, B and C are the inputs of full adder and sum and carry are the outputs of full adder while its schematic is shown in figure 4.

Table 2: Truth Table of Full Adder

A	B	C	Sum	Carry
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

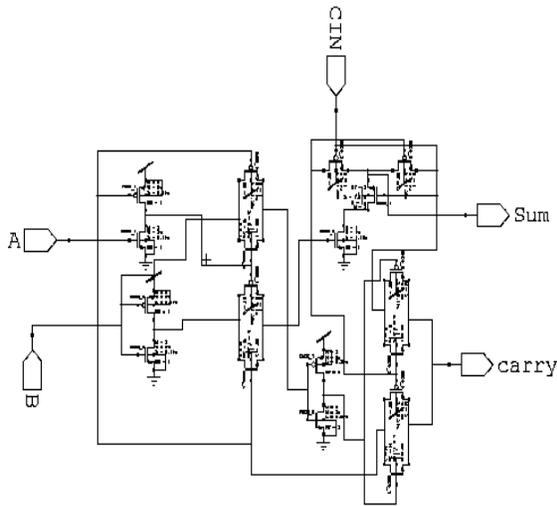


Figure 5:- Schematic of Full Adder

allowing an output port to assume a high impedance state in addition to the 0 and 1 logic levels. This allows multiple circuits to share the same output line or lines. Buffer can be thought of as a switch. If B is on, the switch is closed. If B is off, the switch is open. The operational truth table of tri-state buffer is shown in table where A is the data input and Y (output) schematic is shown in figure 5.

Table 3: Truth Table of Buffer

A (In)	Output
0	0
1	1

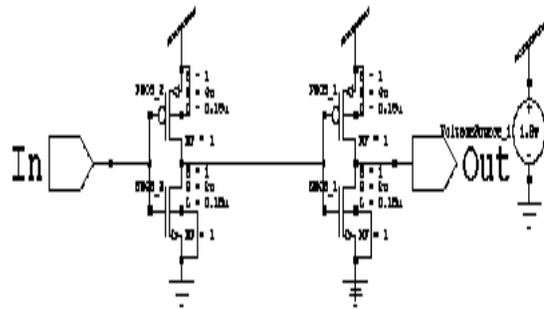


Figure 6: Schematic of Buffer

Schematic of 4 bit ripple carry adder: - The block diagram of 4-bit Ripple Carry Adder are shown in fig 6 In this case there are 4 adders are connected parallel. The full adder is capable of adding only two digit binary numbers along with the carry input. But in practice we need to add binary number which is much longer then just one bit. To add two n-bit binary numbers we need to use the n-bit parallel adder shown in fig 7. It uses the number of full adder in cascade. The carry output of the previous full adder is connected to the carry input to the next full adder as shown in fig 7.

Schematic of Buffer: - Buffer has 2-state logic effectively removes the output from the circuit by

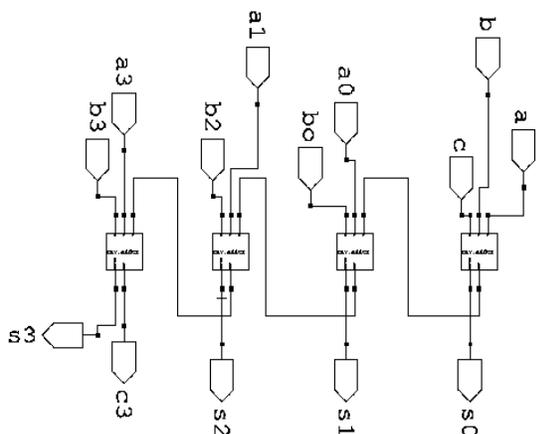


Figure 7: Schematic of 4 bit Ripple Carry Adder

Schematic of NOT: - In digital logic, NOT gate is a logic gate which implements logical negation. An inverter circuit outputs a voltage which represents the opposite logic-level to its input. NOT gate is a logic gate having one input and one output and is also known as inverter. The operational truth table of NOT gate is shown in table while its schematic is shown in figure 7.

Table 4: Truth Table of NOT gate

Input	Output
0	1
1	0

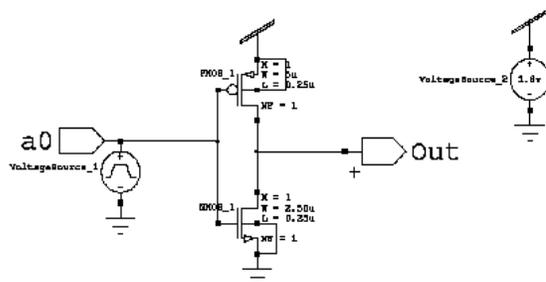


Figure 8:- Schematic of NOT Gate

Schematic of XOR Gate:-The operational truth table of XOR gate is shown in table 4.6 where A, B are the inputs of XOR gate. If two input of XOR gate are low then output will be low and if the two input of XOR gate are high then output is also low. If one input is low and other input is high then output will be high. While its schematic is shown in figure 4.8.

Table 6: Truth Table of XOR gate

A	B	Output
0	0	0
0	1	1
1	0	1
1	1	0

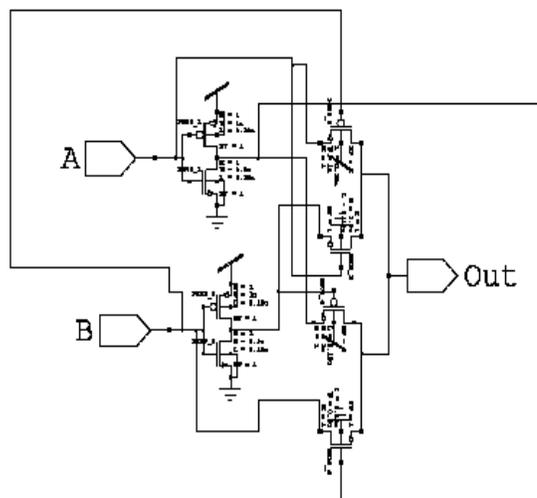


Figure 9: Schematic of XOR gate

Schematic of Half Adder (A + B):-Half adder is a combinational logic circuit with two inputs and two outputs. It is the basic building block for addition of two single bit numbers. Two one-bit binary numbers A and B is added by the half adder. It has two outputs namely “carry” and “sum”. The half-adder design incorporates an XOR gate for S and an AND gate for C. With the addition of an OR gate to combine their carry outputs, two half adders can be combined to make a full adder. The operational truth table of half adder is shown in table where A, B are the inputs and sum and carry are the outputs of half adder while its schematic shown in figure 9.

Table 7: Truth Table of Half Adder

A	B	Sum	Carry
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1

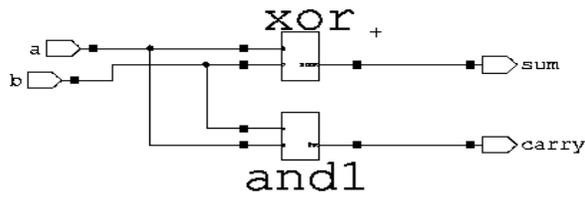


Figure 10: Schematic of half adder

Schematic of 2 – bit Vedic Multiplier

The Schematic of 2 – bit Vedic Multiplier design is shown in Figure 4.12. In the 2 bit Multiplier there are basically 2 half adder is used which is combinations of XOR Gate, AND gate is used and four AND Gate is used. In the 2-bit Vedic multiplier a0, b0, a1, b1 are the input of multiplier and s0, s1, c2, s2 are the output of Vedic Multiplier which is shown in fig:-

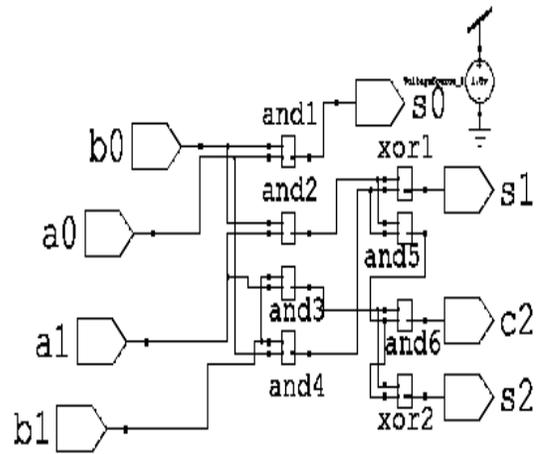


Figure 11: Schematic of 2 bit Vedic Multiplier

The Schematic of 4-bit Vedic Multiplier is shown in figure

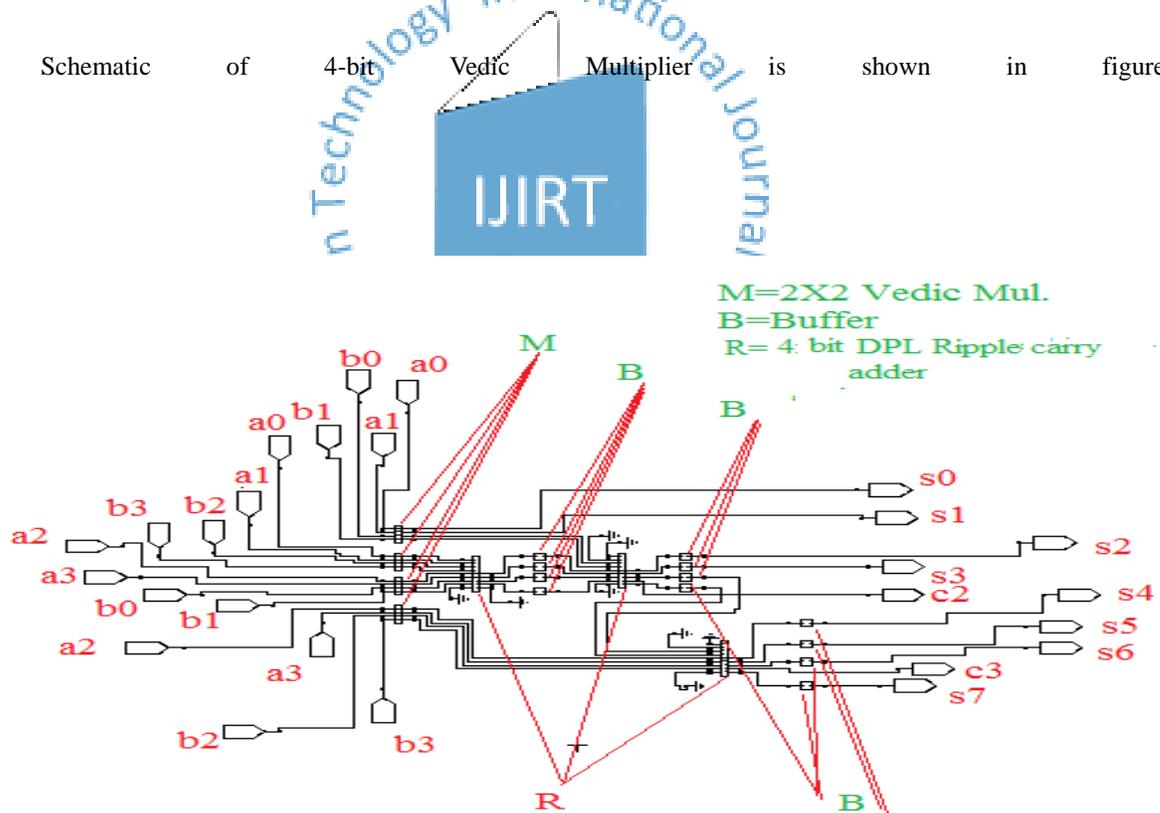


Figure 12: Schematic of 4 – bit Vedic Multiplier using double pass-transistor Logic (DPL) Adder

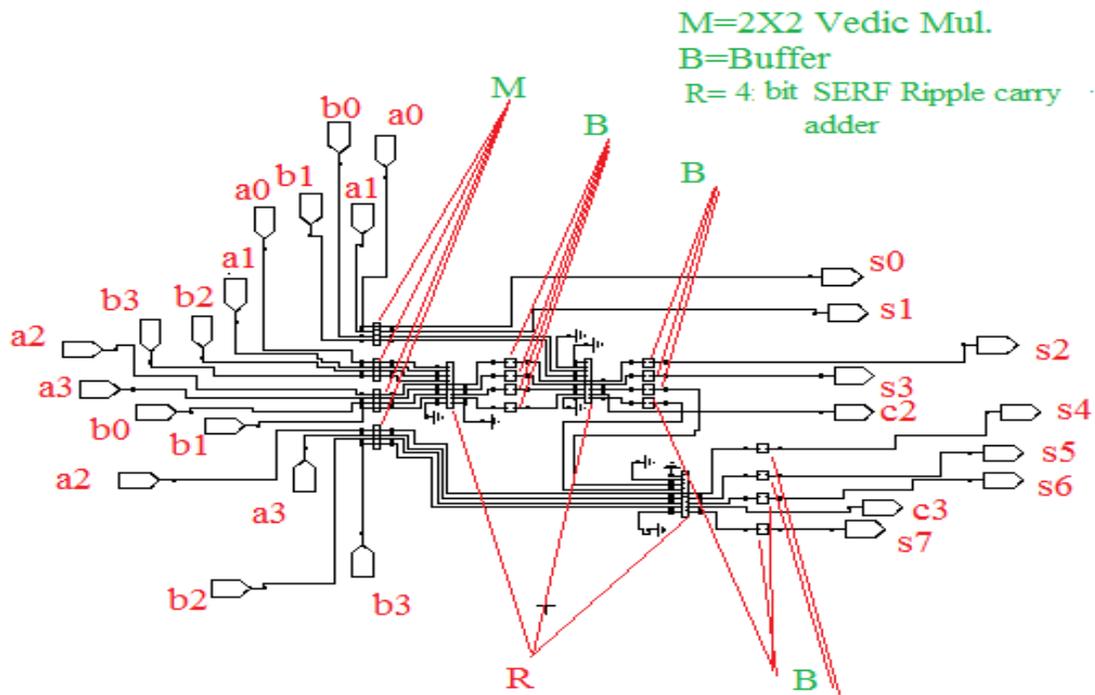


Figure 13: Schematic of 4 – bit Vedic Multiplier using SERF Adder

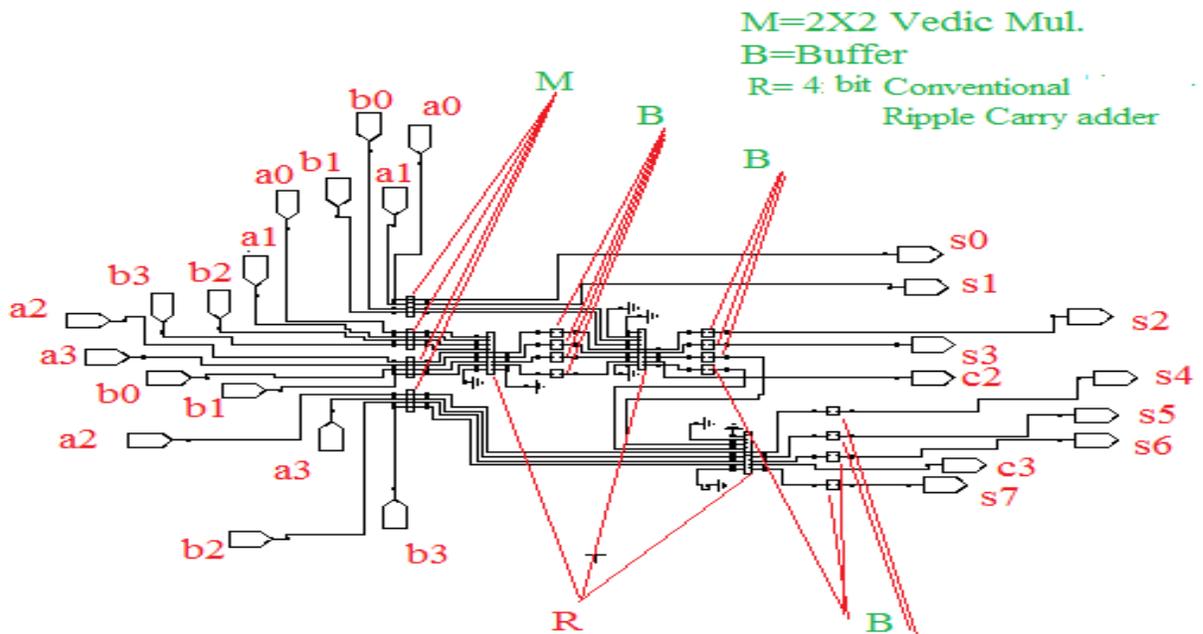
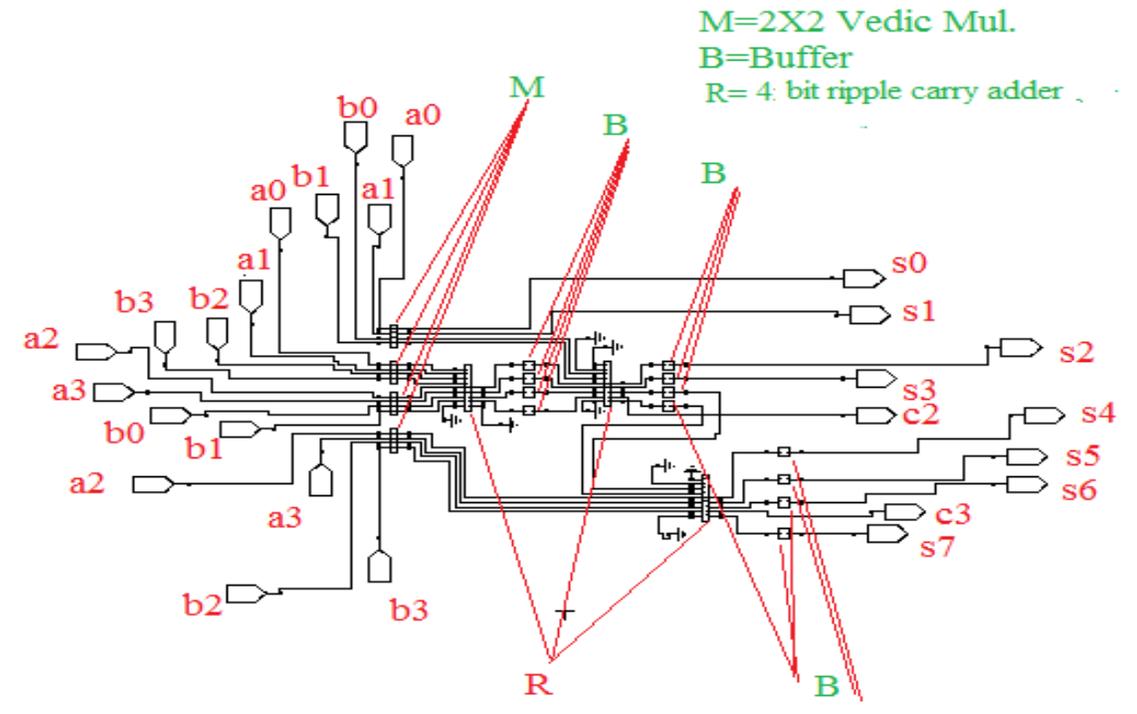
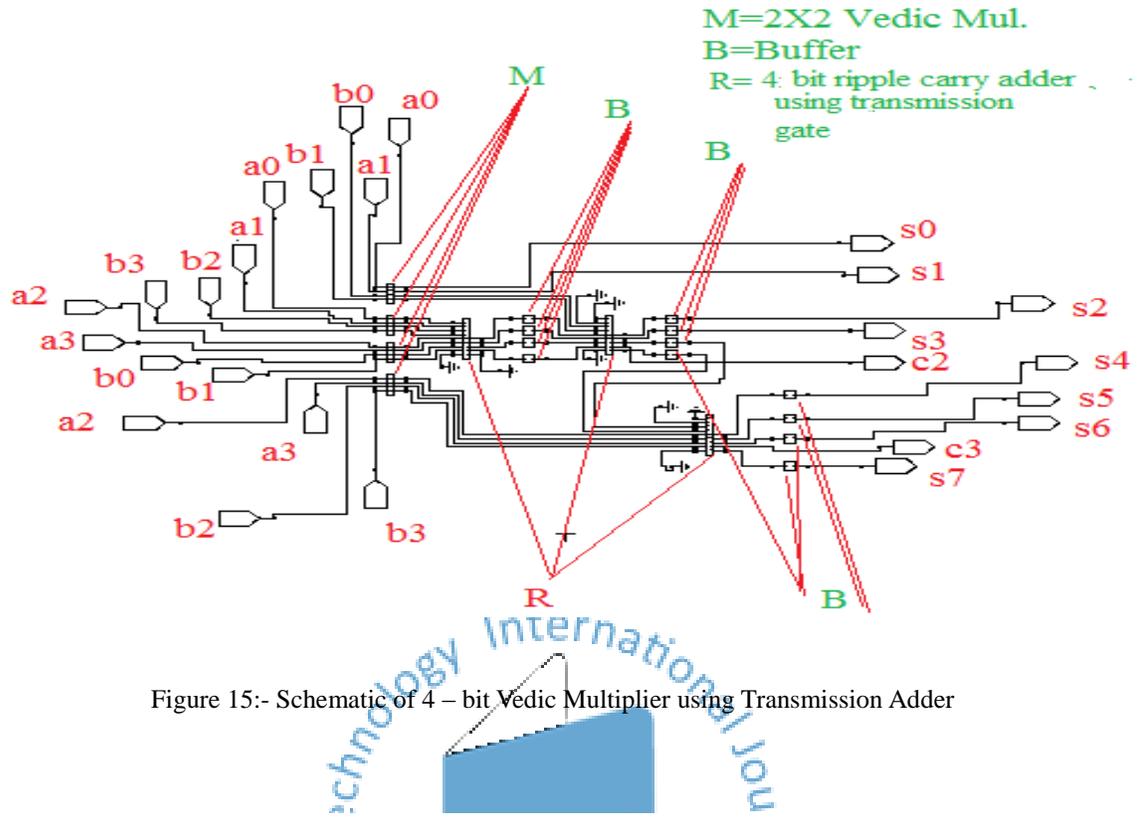


Figure 14: Schematic of 4 – bit Vedic Multiplier using Conventional Adder



Simulation Waveform of Full Adder

The v (a), v (b), v (cin), are the input of full adder and v (sout), v (cout) are the output of Full Adder. Where the v (sout) indicate the sum and v (cout) indicate the carry. The simulation waveform of full adder is shown in Figure 17.

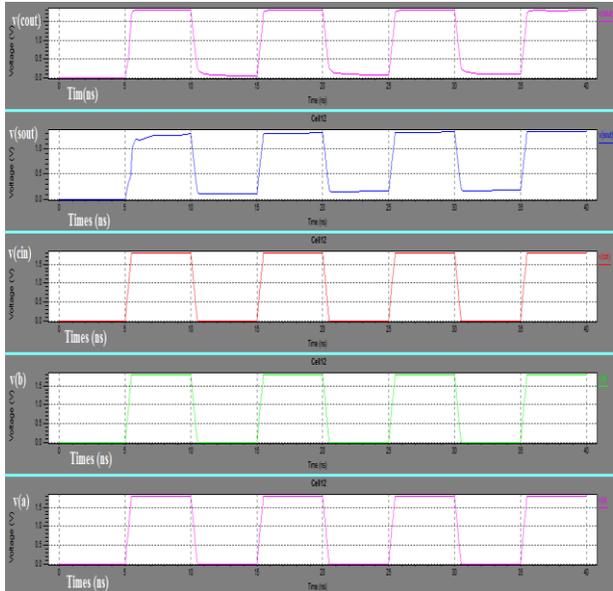


Figure 17: Simulation Waveform of SERF Adder

Simulation Waveform of Buffer

The v (In) is the input of buffer and v (out) are the output of Buffer. The simulation waveform of buffer is shown in Figure 18.

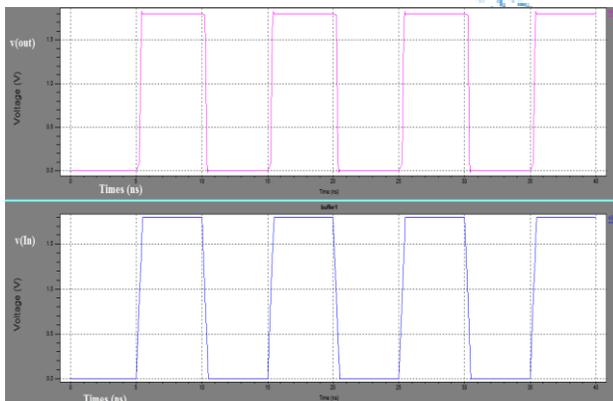


Figure 18: Simulation Waveform of buffer

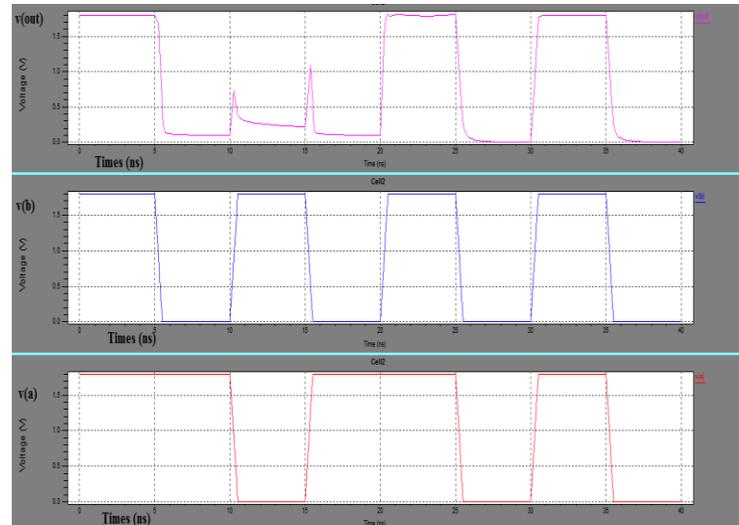


Figure 19: Simulation Waveform of AND Gate

Simulation Waveform of XOR Gate

The v (A), v (B) is the input of XOR gate. The v(out) is the output of XOR gate. The simulation waveform of XOR gate is shown in Figure 20

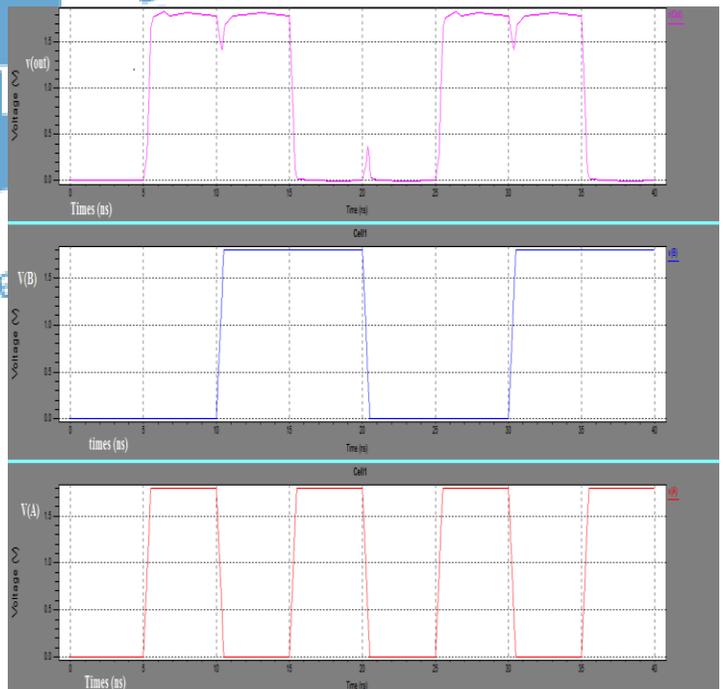


Figure 20: Simulation Waveform of XOR gate

Simulations Waveform of Ripple Carry Adder

The v (a), v (b), v(c) is the input of Ripple Carry Adder and v (s1) and v(c) are the output of 4 bit Ripple Carry Adder. The v(c) is the carry which will be generated at the end of Ripple Carry Adder and v (s1) will work as the sum of 4-bit Ripple Carry Adder. The simulation waveform of Ripple Carry Adder is shown in Figure

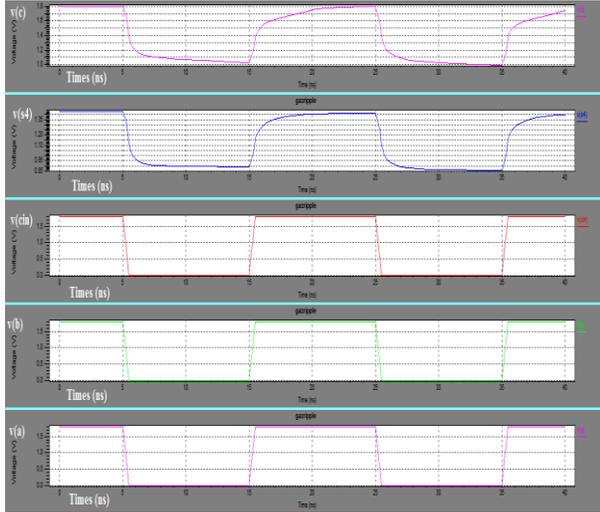


Figure 21:-Simulations Waveform of Ripple Carry Adder

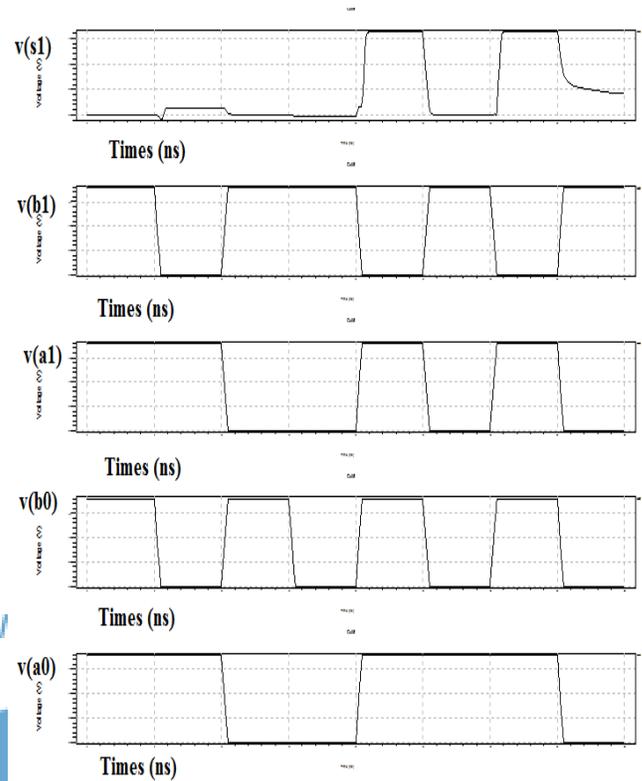


Figure 23: Simulation Waveform of S1

Simulations Waveform for Vedic Multiplier:- In the 2 bit Vedic multiplier the v (a0) ,v (a1)and v(b0),v (b1) are the input of 2-bit Vedic multiplier and v (s0), v(s1), v(s2), v(c2) is the output of 2-bit Vedic multiplier. Where v (c2) is indicates the carry. The simulations waveform is shown follows:-

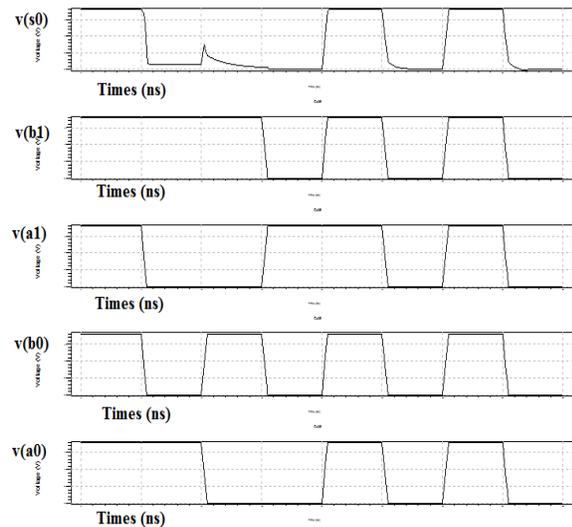


Figure 22: Simulation Waveform of S0

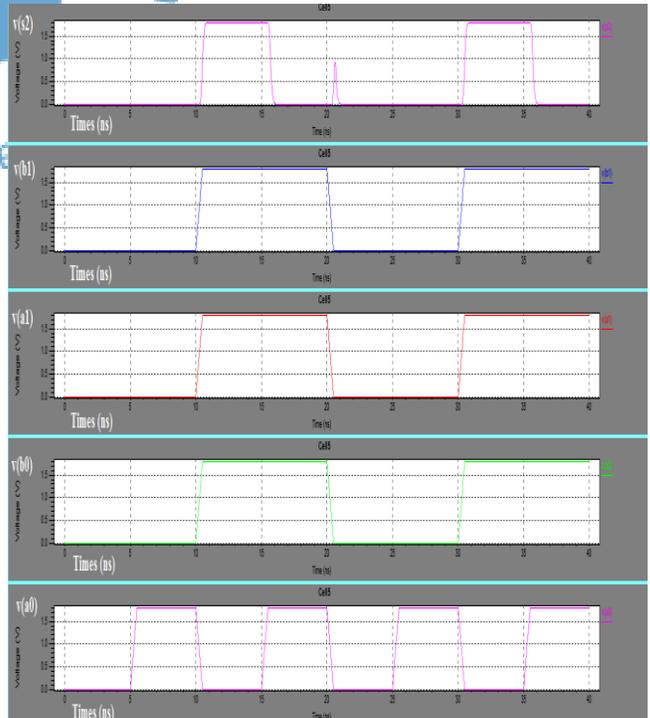


Figure 24:-Simulation Waveform of S2

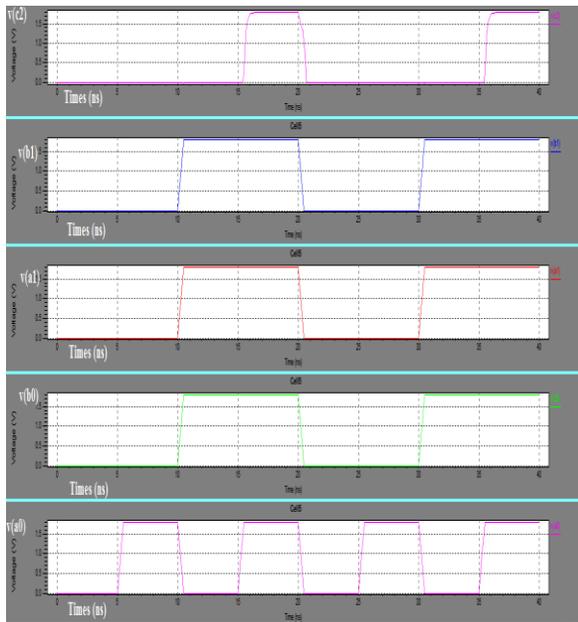


Figure 25: Simulation Waveform of C2

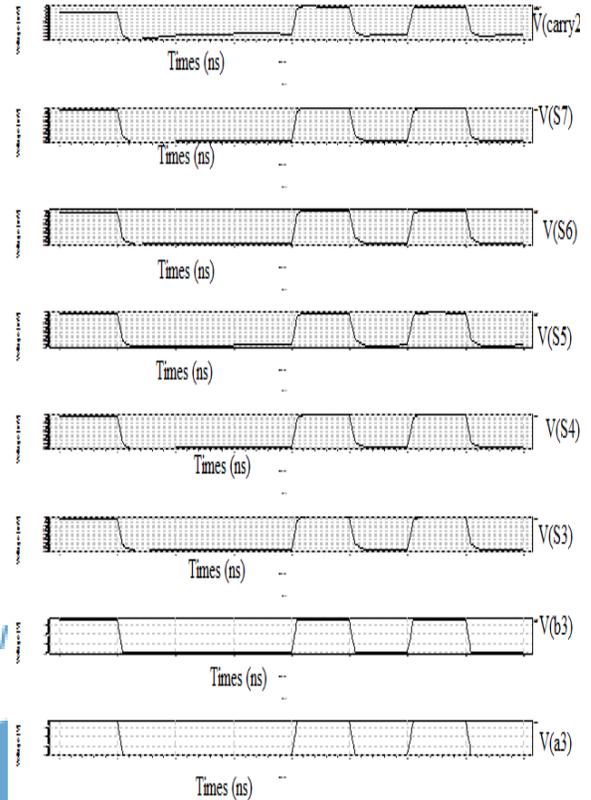


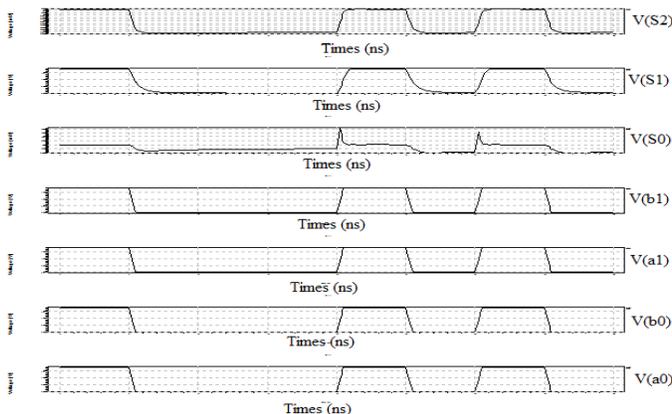
Figure 26: Simulation Waveform of 4 bit Vedic Multiplier Using SERF Adder

Figure 27: Simulation Waveform of 4 bit Vedic Multiplier Using SERF Adder

IV. RESULTS OF VEDIC MULTIPLIERS

Table 8: Number of transistor of 4-Vedic bit Multipliers

Sr. No.	Multiplier Name	No. of Transistor used
1.	2 X 2 Bit Multiplier	58
2.	4 x 4 bit multiplier using SERF Adder	400
3.	4 x 4 bit multiplier using conventional Adder	520
4.	4 x 4 bit multiplier using 1-bit Full Adder	520



5.	4 x 4 bit multiplier using DPL Adder	616
6.	4 x 4 bit multiplier using Transmissions gate Adder	445

6.	4 x 4 bit multiplier using Transmissions gate Adder	0.3.11µw
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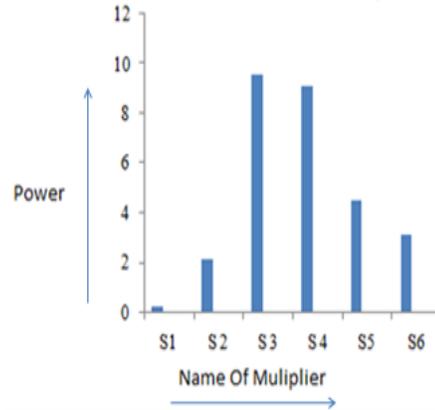
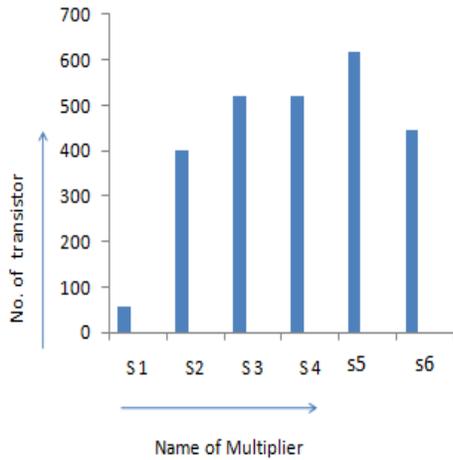


Figure 28:- No. of transistor vs. Name of Multiplier

Fig 29:-Average Power vs. Name of Multiplier

S1= 2x2 Vedic Multiplier.

S2 = 4 x 4 bit multiplier using SERF Adder. S3=4 x 4 bit multiplier using conventional Adder.

S4=4 x 4 bit multiplier using 1-bit Full Adder. S5=4 x 4 bit multiplier using DPL Adder.

S6=4 x 4 bit multiplier using Transmissions gate Adder.

S1= 2x2 Vedic-Multiplier.

S2 = 4 x 4 bit multiplier using SERF Adder.

S3= 4 x 4 bit multiplier using conventional Adder.

S4= 4 x 4 bit multiplier using 1-bit full adder

S5= 4 x 4 bit multiplier using DPL Adder.

S6=4 x4 bit multiplier using transmission gate adder.

Table 9: Power dissipations of 4-Vedic bit Multipliers

Table 10: Propagations Delay of 4-Vedic bit Multiplier

Sr. No.	Multiplier Name	Power Dissipations At 500 MHz
1.	2 X 2 Bit Multiplier	0.21µw
2.	4 x 4 bit multiplier using SERF Adder	2.11µw
3.	4 x 4 bit multiplier using conventional Adder	9.55µw
4.	4 x 4 bit multiplier using 1-bit Full Adder	9.07µw
5.	4 x 4 bit multiplier using DPL Adder	4.470µw

Sr. No.	Multiplier Name	Propagations Delay
1.	2 X 2 Bit Multiplier	0.09ns
2.	4 x 4 bit multiplier using SERF Adder	1.23ns
3.	4 x 4 bit multiplier using conventional Adder	0.39ns
4.	4 x 4 bit multiplier using 1-bit Full Adder	1.45 ns
5.	4 x 4 bit multiplier using DPL Adder	2.31ns
6.	4 x 4 bit multiplier using Transmissions gate Adder	3.01 ns

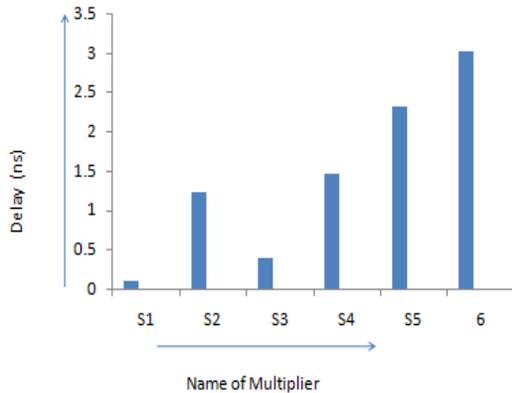


Figure 30:- Delay vs. Name of Multiplier

S1= 2x2 Vedic Multiplier.

S2 = 4 x 4 bit multiplier using SERF Adder.

S3= 4 x 4 bit multiplier using conventional Adder.

S4= 4 x 4 bit multiplier using 1-bit full adder

S5= 4 x 4 bit multiplier using DPL Adder.

S6=4 x4 bit multiplier using transmission gate adder.

V. CONCLUSION

This work focuses on optimization of power and high performance of CMOS Vedic multiplier. The comparative analyses of Low power Vedic Multiplier using different adder logic style are done in this thesis. Result shows that the power consumed by the 2 bit Vedic multiplier is $0.21\mu\text{w}$ and the delay is 0.09ns . The average power consumed by the 4 bit Vedic multiplier using SERF Adder is $2.11\mu\text{w}$ and the delay 1.23ns . The power consumed by the 4 bit Vedic multiplier by using DPL Adder is $4.470\mu\text{w}$ and the 4 bit Vedic multiplier using DPL Adder is 2.31ns . The power consumed by the 4 bit Vedic multiplier by using transmission gate is $3.11\mu\text{w}$ and the delay is 3.01ns . The power consumed by the 4 bit Vedic multiplier using Conventional Adder is $9.55\mu\text{w}$ and the delay is 0.39ns . The power consumed by the 4 bit Vedic multiplier using 1-bit full Adder is $9.07\mu\text{w}$ and the delay is 1.45ns . Vedic Multiplier when used by using SERF adder is having lesser power dissipation.

VI. FUTURE SCOPE

Only one aspect of Vedic mathematics is touched in this work but there are so many other algorithms can be explored for multiplication. There is further scope

of improve implementation of the multiplier presented in this thesis to optimize the area & static power dissipations. Multiplication and divisions have innumerable application which can be implemented using Vedic mathematics. Floating point Vedic processor could be the extensions of this work.

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